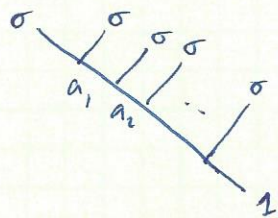


• Ising anyons  $\sigma \cdot \sigma = 1 \neq \tau$



$$N_{\sigma\sigma}^{a_1}, N_{\sigma\sigma}^{a_2}, N_{\sigma\sigma}^{a_3}, \dots, N_{\sigma\sigma}^{a_{n-1}}$$

$$N_{a\sigma}^b \rightarrow (N_\sigma)_a^b = \begin{pmatrix} 1 & \sigma & \tau \\ 0 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 1 & 0 \end{pmatrix}$$

eigenvalues  $\lambda = 0, \pm\sqrt{2}$

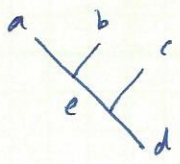
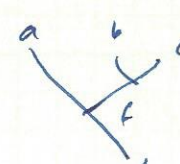
$\therefore$  # of states  $\sim (\sqrt{2})^n$

"quantum dimension"  $d_\sigma = \sqrt{2}$   
# states per particle

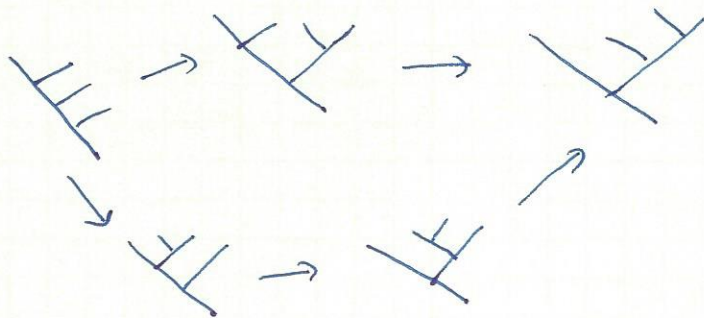
• Quantum dim : Total q.d.  $D \equiv \sqrt{d_a^2}$

generalizes |D| discr. gp. of a lattice.

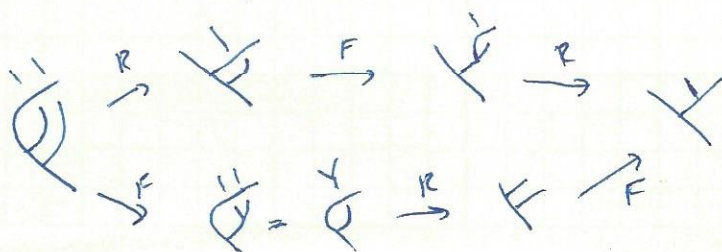
$$S = \frac{1}{\sqrt{D}} \sum_c d_c (R_c^{ab} R_c^{ba}) = \mathcal{O}$$

•  =  $\sum_f \mathcal{O} [F_{abc}^d]_e^f$  

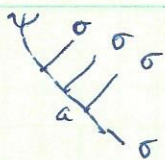
• Pentagon:



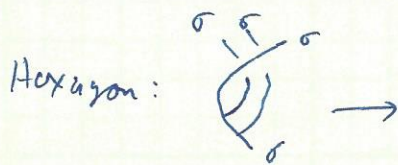
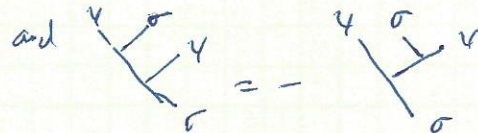
• Hexagon:



Ising anyon: Pentagon:



$f_{\sigma\sigma\sigma} = \frac{1}{\sqrt{2}} \begin{pmatrix} 1 & -1 \\ 1 & 1 \end{pmatrix}$



$R_{\sigma\sigma}^{\sigma} = e^{-\pi i/8}$ ,  $R_{\sigma\sigma}^{\sigma} = e^{3\pi i/8}$   
 $R_{\sigma\sigma}^{\sigma} = -i = e^{\pi i/8} e^{3\pi i/4}$

$\theta_{\sigma} = e^{+\pi i/8}$ ,  $\theta_4 = -1$ ,  $\theta_1 = 1$ ,  $\theta_a = \pm (R_{\sigma\sigma}^a)^*$

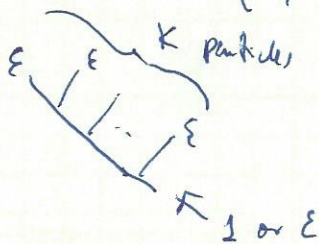
~~Electron~~

$\frac{1}{\sqrt{D}} \sum d_a^2 \theta_a = e^{\pi i/8} \Rightarrow c = 1/2$   $|\rho = \theta_a|$

Fibonacci anyon 1,  $\epsilon$

$\epsilon \cdot \epsilon = 1 + \epsilon$

$N_{\epsilon} = \begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix} \Rightarrow d_{\epsilon} = \frac{1+\sqrt{5}}{2} = \phi$  "golden ratio"



$n_k^1 = n_{k-1}^{\epsilon}$

$n_k^{\epsilon} = n_{k-1}^{\epsilon} + n_{k-1}^1 = n_{k-1}^{\epsilon} + n_{k-2}^{\epsilon}$

Fibonacci numbers

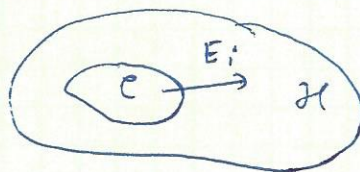
$F_{\epsilon\epsilon\epsilon}^{\epsilon} = \begin{pmatrix} \phi^{-1} & \phi^{-1/2} \\ \phi^{-1/2} & \phi^{-1} \end{pmatrix}$

$R_{\epsilon\epsilon}^{\epsilon} = e^{4\pi i/5}$

$R_{\epsilon\epsilon}^{\epsilon} = e^{-3\pi i/5}$

• Error-correcting Code

$$P_c E_j^\dagger E_i P_c = c_{ij} P_c$$



$C \subset \mathcal{H}$   
 $\uparrow$   
 code subspace       $\nwarrow$   
                                  Hilbert space

e.g. Show 9-qubit code:  $|0\rangle_L = \frac{1}{\sqrt{2}}(1000\rangle + |1111\rangle) \frac{1}{\sqrt{2}}(1000\rangle + |1111\rangle) \cdot \frac{1}{\sqrt{2}}(1000\rangle + |1111\rangle)$

$$|1\rangle_L = \frac{1}{\sqrt{2}}(1000\rangle - |1111\rangle) \frac{1}{\sqrt{2}}(1000\rangle - |1111\rangle) \cdot \frac{1}{\sqrt{2}}(1000\rangle - |1111\rangle)$$

Compare to def. of a top. phase:

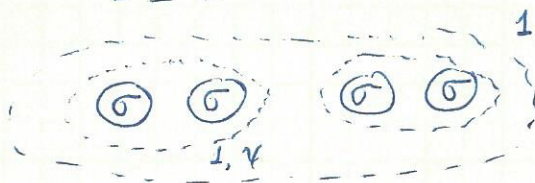
$$\langle a | x | b \rangle = c \delta_{ab} + O(e^{-L/\xi})$$

$\uparrow$   
 The different ground states of a top. phase (e.g. on an n-punctured plane) can serve as a code subspace

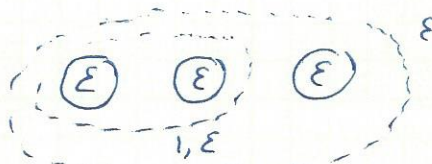
• Ising anyons: Dense encoding



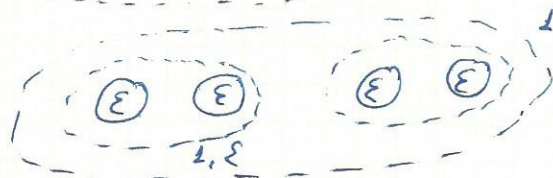
Sparse encoding



• Fibonacci anyons: Dense encoding

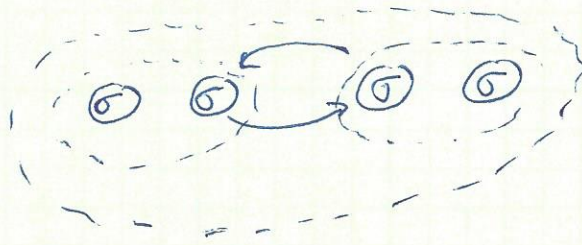


Sparse encoding



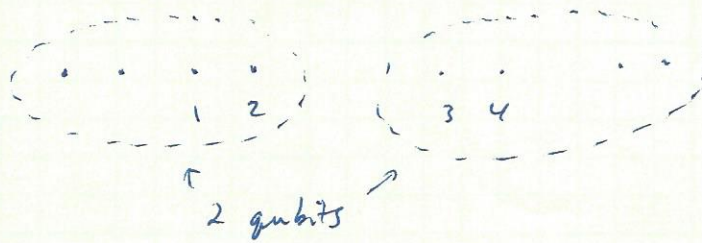
Gate Sets:  $H, \Lambda(\sigma_z), T$  form a complete set of universal gates

Ising  $H$ :



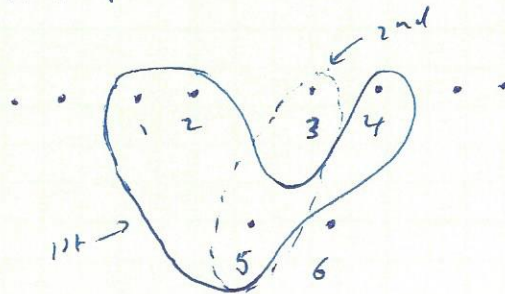
$$\Lambda(\sigma_z) = e^{i\frac{\pi}{4}} e^{-i\frac{\pi}{4}\sigma_1\sigma_2\sigma_3\sigma_4} e^{-\frac{\pi}{4}\sigma_1\sigma_2} e^{-\frac{\pi}{4}\sigma_3\sigma_4}$$

where



Need to apply  $e^{-i\frac{\pi}{4}\sigma_1\sigma_2\sigma_3\sigma_4}$

Use an ancilla and perform measurements:

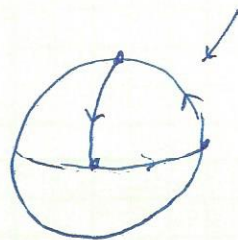


Bravyi '05

$$\begin{aligned} e^{i\frac{\pi}{4}\sigma_1\sigma_2\sigma_3\sigma_4} &= 2 e^{\frac{\pi}{4}\sigma_3\sigma_6} \Pi_+^{(2)} \Pi_-^{(4)} \\ &= 2i e^{\frac{\pi}{2}\sigma_1\sigma_2} e^{\frac{\pi}{2}\sigma_3\sigma_4} e^{\frac{\pi}{4}\sigma_3\sigma_6} \Pi_-^{(2)} \Pi_-^{(4)} \end{aligned}$$

$$T = \begin{pmatrix} 1 & \\ & e^{i\pi/4} \end{pmatrix}$$

$$S = T^2 = \begin{pmatrix} 1 & \\ & i \end{pmatrix} \quad \text{---} \quad \text{---} \quad \text{---} \quad \text{---}$$



octant on Bloch sphere

Need to produce  $\frac{1}{\sqrt{2}}(|0\rangle + e^{i\pi/4}|1\rangle)$

Then  $T(a|0\rangle + b|1\rangle)$  can be obtained from

$$(a|0\rangle + b|1\rangle) \frac{1}{\sqrt{2}}(|0\rangle + e^{i\pi/4}|1\rangle)$$

↓ Joint parity

$$a|0\rangle|0\rangle + b e^{i\pi/4}|1\rangle|1\rangle$$

↓ CNOT

$$(a|0\rangle + b e^{i\pi/4}|1\rangle)|1\rangle$$

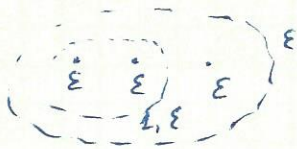
15 copies of  $\frac{1}{\sqrt{2}}(|0\rangle + e^{i\pi/4}|1\rangle)$   $\xrightarrow{\text{Reed-Muller code}}$   $\frac{1}{\sqrt{2}}(|0\rangle_L + e^{i\pi/4}|0\rangle_L)$

error  $\epsilon$

1 copy of  $\frac{1}{\sqrt{2}}(|0\rangle + e^{i\pi/4}|1\rangle)$   
w/ error  $35\epsilon^3$

How to make T?

• Fibonacci



universal gate set!

$$\begin{matrix} \text{CNOT} & = & \begin{pmatrix} e^{4\pi i/5} & \\ & e^{-3\pi i/5} \end{pmatrix} \\ \text{CNOT} & = & \begin{pmatrix} \phi^{-1} & \phi^{-1/2} \\ \phi^{-1/2} & \phi^{-1} \end{pmatrix} \begin{pmatrix} e^{4\pi i/5} & \\ & e^{-3\pi i/5} \end{pmatrix} \end{matrix}$$

$$\begin{pmatrix} \phi^{-1} & \phi^{-1/2} \\ \phi^{-1/2} & \phi^{-1} \end{pmatrix}$$

How to obtain Fibonacci anyons?

Algorithms tailored to Fibonacci!

Need  $|j\rangle \rightarrow \frac{1}{\sqrt{2^n}} \sum_{k=0}^{2^n-1} e^{2\pi i jk/2^n} |k\rangle$

Quantum Fourier Transform

have  $e^{i\pi}$ ,  $e^{2\pi i/4}$ ,  $e^{2\pi i/5}$ ,  $e^{2\pi i/10}$ ,  $e^{2\pi i/20}$