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Tensor Networks - A natural way to represent
~~both an analytic and a numerical~~
quantum states when there is low entanglement
(normally the case): MPS, PEPS, MERA

- 1D Matrix Product States (well understood)
time evolving ground states
- higher dimension
- end of second lecture: slides showing some results
for spin liquids

References

U Schollwock arxiv 1008.3477

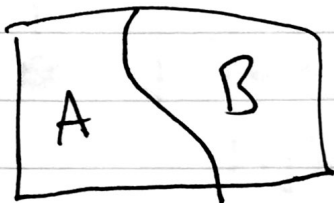
Software: 1) itensor.org

2) juliaLang.org - very nice new language,
osp. for tensor
networks

Web site

<https://eee.uci.edu/165/48480>

Review: Schmidt Decomposition



$$|\psi\rangle = \sum_{a,b} \psi_{ab} |a\rangle |b\rangle$$

Schmidt Form: $|\psi\rangle = \sum_{\alpha} \lambda_{\alpha} |\alpha\rangle_A |\alpha\rangle_B$

- diagonal
- Always possible

- Calculation: (1) SVD: $\psi_{ab} = [U D V^T]_{ab}$

$$D = \begin{pmatrix} \lambda_1 & & 0 \\ & \lambda_2 & \\ 0 & & \dots \end{pmatrix} \quad U, V \text{ unitary} \quad |\alpha\rangle_A = \sum_a U_{a\alpha} |a\rangle$$

(2) Reduced density matrix

$$\rho_{aa'}^A = \sum_b \psi_{ab} \psi_{a'b}^* = [U D V^T V D^T U^T]_{aa'} \\ = [U D^2 U^T]_{aa'}$$

$\lambda_{\alpha}^2 =$ eigenvalue of ρ^A and ρ^B
eigenvectors of ρ^A form columns of U , etc.

Von Neumann entanglement entropy

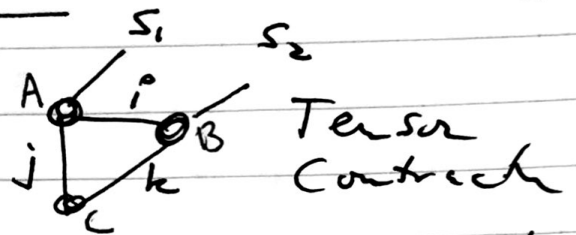
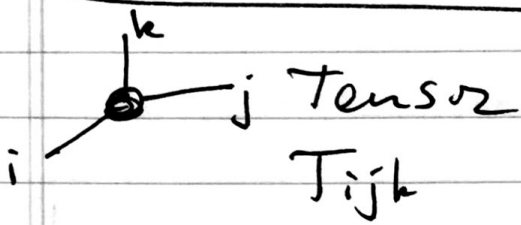
$$S = \sum_{\alpha} -\lambda_{\alpha}^2 \ln \lambda_{\alpha}^2 \quad \lambda_{\alpha}^2 = \text{prod of } \alpha \text{ - pair of states}$$

For ground states

$$S \sim \text{Area-Law}$$

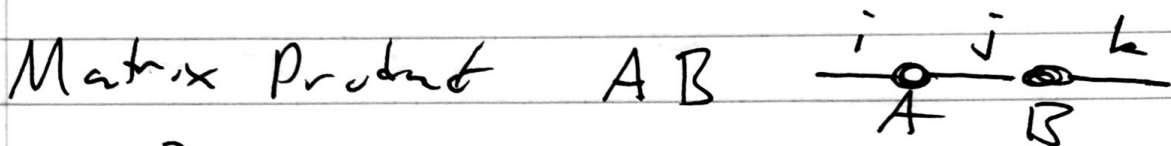
Tensor Network Diagrams

A_{jis_1} B_{is_2k}
 C_{jk}

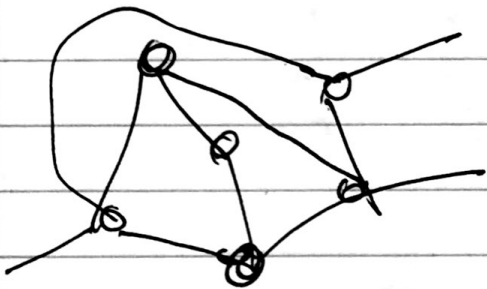


3 tensors. Internal ~~legs~~ indices summed over. External indices define the final tensor result.

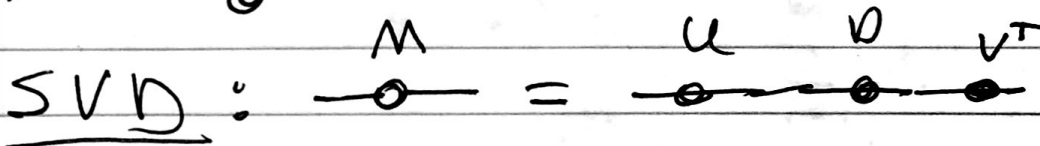
~~Matrix~~ Vector \bullet $\text{---} i$ Matrix $\text{---} \bullet$



$$[AB]_{ik} = \sum_j A_{ij} B_{jk}$$



Tensor networks can get complicated.



$$M = U D V^T$$

Truncating Low Probability States

L	R
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Suppose we have λ_i $\sum \lambda_i^2 = 1$ $S=1$ char

We want to throw
away states with
small λ_i

(Schmidt states)

$$\text{Error} = \sum_{i=1}^m \lambda_i |i\rangle_L |i\rangle_R \rightarrow \sum_{i=1}^m \lambda_i |i\rangle_L |i\rangle_R$$

Error \approx $|i\rangle_L$ $\tilde{\psi}$

$$(\langle \psi | - \langle \tilde{\psi} |) (\psi \rangle - |\tilde{\psi} \rangle)$$

$$\langle \psi | \tilde{\psi} \rangle = \sum_{i=1}^{2^{N/2}} \sum_{i'=1}^m \lambda_i \lambda_{i'} \underbrace{\langle i | i' \rangle}_{\delta_{ii'}} = \sum_{i=1}^m \lambda_i^2$$

$$\delta_{ii'} = \langle \tilde{\psi} | \tilde{\psi} \rangle$$

$$= \langle \psi | \psi \rangle - \langle \tilde{\psi} | \tilde{\psi} \rangle = 1 - \sum_{i=1}^m \lambda_i^2 = \sum_{i=m+1}^{2^{N/2}} \lambda_i^2$$

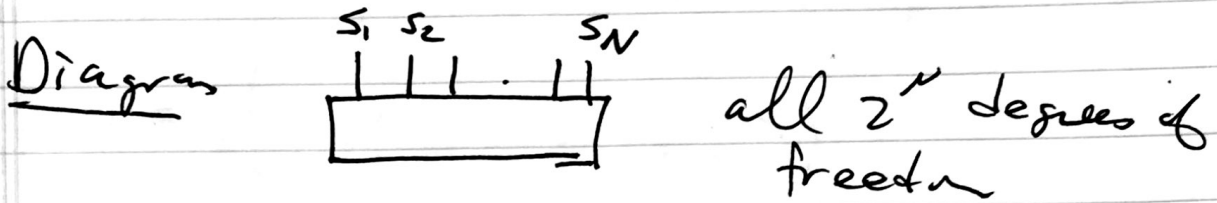
error $\approx \psi$ is $\sum_{i=m+1}^{2^{N/2}} \lambda_i^2 =$ "truncation error"
or "discarded weight"

This is ~~also~~ the error on the
density matrix, roughly.

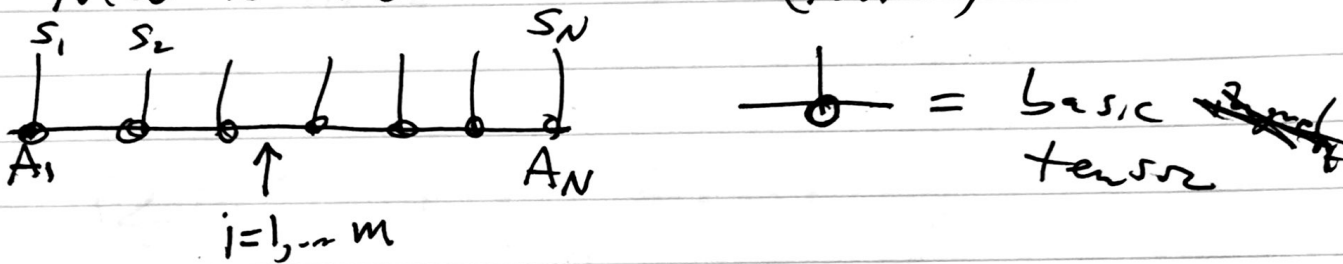
This is the key TN approximation: a
Low entanglement approximation

Suppose we had a wavefunction from exact diagonalization of a spin chain

$$\Psi(s_1, s_2, \dots, s_N) \quad 2^N \text{ numbers} \quad (N \leq \sim 50)$$



As an ansatz, let's propose a Matrix Product State (MPS) is a TN:



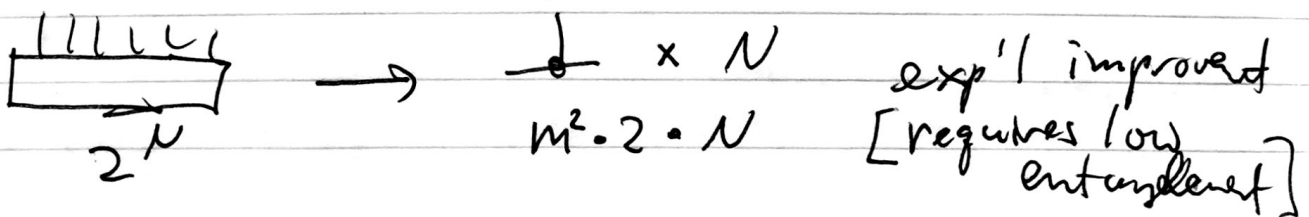
Algebraically $\Psi(s_1, \dots, s_N) = \overset{\text{row vec}}{\rightarrow} A^1[s_1] \overset{\text{col vec}}{\leftarrow} A^2[s_2] \dots A^N[s_N]$

Regard $A^i[s_i]$ as two matrices

$$A^i[\uparrow]_{ij} = \downarrow \quad \text{and} \quad A^i[\downarrow]_{ij} = \uparrow$$

Given s_1, \dots, s_N , pick which matrix to use on each site, take product.

This is a huge compression



Matrix Product State (MPS)

$$\psi(s_1, \dots, s_N) = \vec{A}_1[s_1] \vec{A}_2[s_2] \dots \vec{A}_N[s_N]$$

Very Simple example: 2 sites, spin 1/2

$$|\psi\rangle = \frac{1}{\sqrt{2}} (|\uparrow\downarrow\rangle - |\downarrow\uparrow\rangle)$$

$$= \sum_{s_1, s_2} \vec{A}_1[s_1] \cdot \vec{A}_2[s_2] |s_1\rangle |s_2\rangle$$

$$= \underbrace{\left(\sum_{s_1} \vec{A}_1[s_1] |s_1\rangle \right)}_{\vec{A}_1} \cdot \underbrace{\left(\sum_{s_2} \vec{A}_2[s_2] |s_2\rangle \right)}_{\vec{A}_2}$$

Let $\vec{A}_1 = \begin{pmatrix} \frac{1}{\sqrt{2}} |\uparrow\rangle \\ -\frac{1}{\sqrt{2}} |\downarrow\rangle \end{pmatrix}$ $\vec{A}_2 = \begin{pmatrix} |\downarrow\rangle \\ |\uparrow\rangle \end{pmatrix}$

By inspection, $|\psi\rangle = \vec{A}_1 \cdot \vec{A}_2$

Three sites: Let $|\psi\rangle = \frac{1}{\sqrt{3}} (|\uparrow\uparrow\downarrow\rangle + |\uparrow\downarrow\uparrow\rangle + |\downarrow\uparrow\uparrow\rangle)$

$$\vec{A}_1 = \begin{pmatrix} |\uparrow\rangle \\ |\downarrow\rangle \end{pmatrix} \quad \vec{A}_3 = \begin{pmatrix} |\uparrow\rangle \\ |\downarrow\rangle \end{pmatrix} \quad \text{then } \vec{A}_2 = \frac{1}{\sqrt{3}} \begin{pmatrix} |\downarrow\rangle & |\uparrow\rangle \\ |\uparrow\rangle & 0 \end{pmatrix}$$

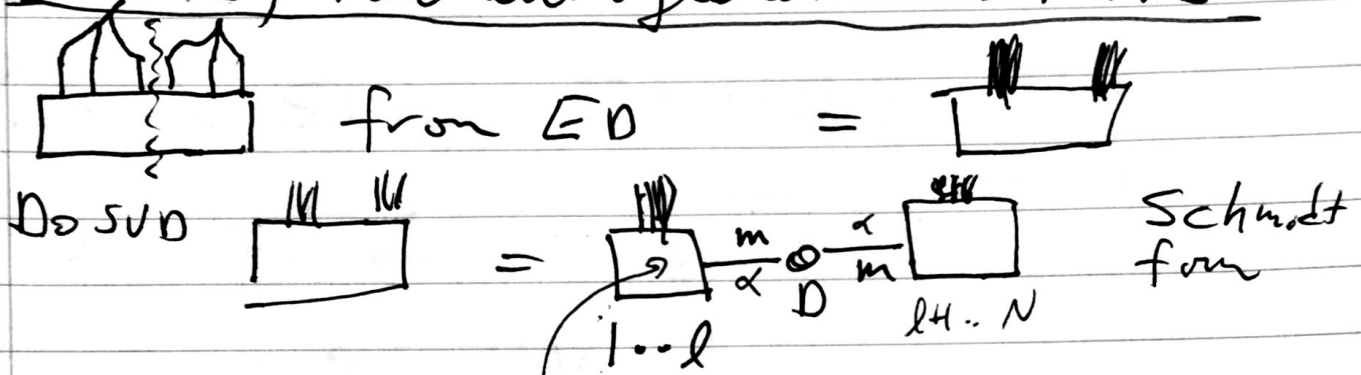
(use \vec{A}_i^T)

More sites: Same structure, bigger matrices, much compressible for low entanglement

6 Builder

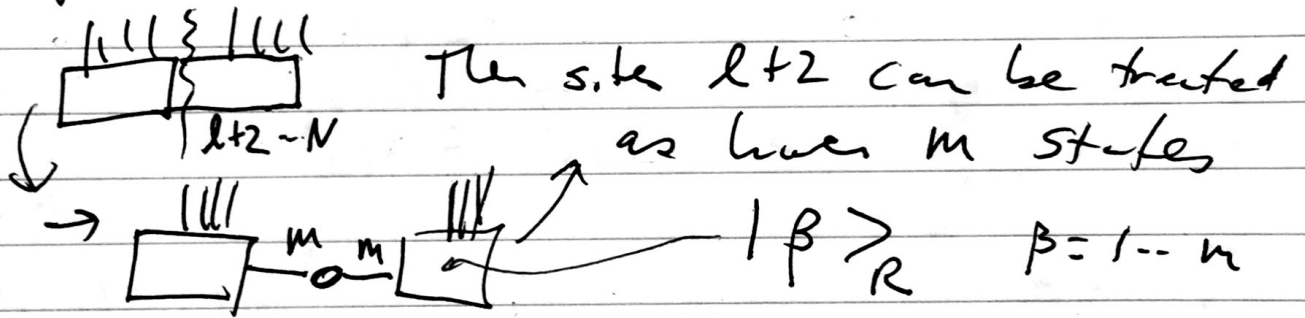
In 1D, low entanglement \Rightarrow MPS

Treat as one big index



Sites $1-l$ can be treated as having only m states $|\alpha\rangle$ $\alpha=1..l$

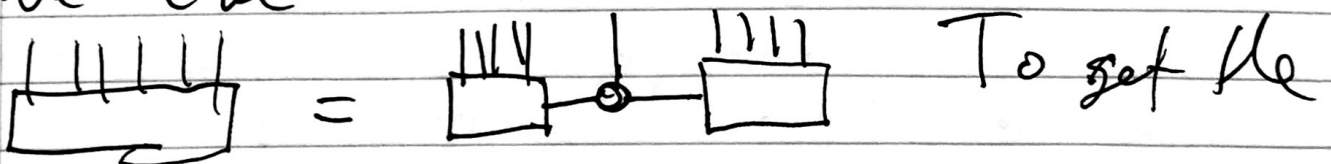
Repeat at one site over



The $|\psi\rangle = \sum_{\alpha=1}^m \sum_{\beta=1}^m \sum_{S_{l+1}=\uparrow, \downarrow} \psi(\alpha, S_{l+1}, \beta)$

$\psi(\alpha, S_{l+1}, \beta) \equiv A[S_{l+1}]_{\alpha\beta} |\alpha\rangle_l |S_{l+1}\rangle | \beta \rangle_r$

We have



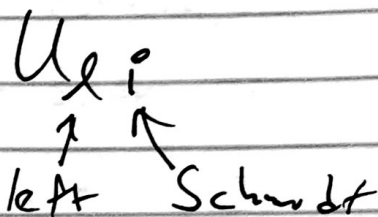
full MPS, you just need to repeat on the other sites.

Orthogonality: crucial for effectively using MPS \downarrow drop t , redefine V

In an SVD, $M = U D V$, $U + V$ are unitary
are row/col. unitary

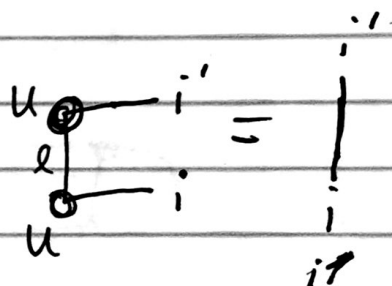
$$U^T U = 1 \quad V V^T = 1$$

Let U have labels indices

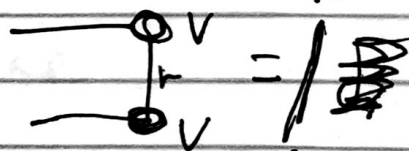


~~V_{lr}~~ V_{lr} \leftarrow right

$$\sum_l U_{li} U_{li}^* = \delta_{ii'} \quad \text{or}$$

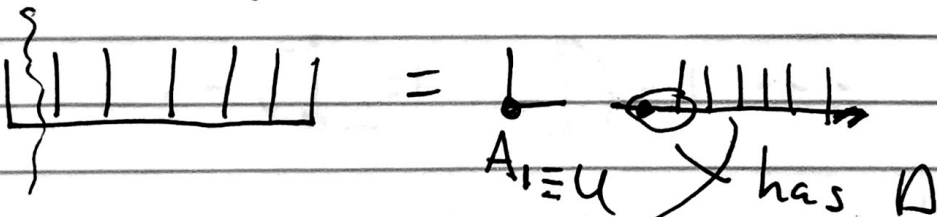


$$\sum_r V_{lr} V_{lr}^* = \delta_{ii'} \quad \text{or}$$



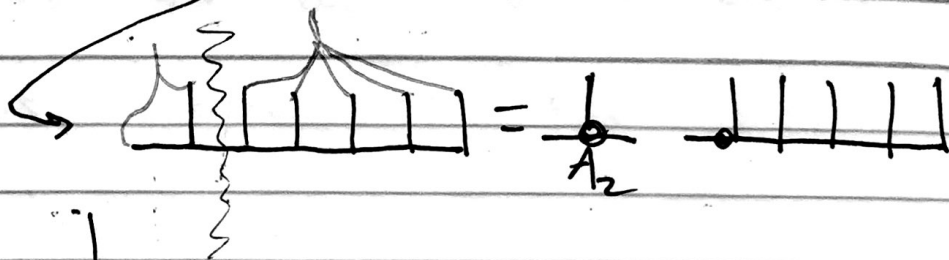
Orthogonality

Converting a Wfn to an MPS

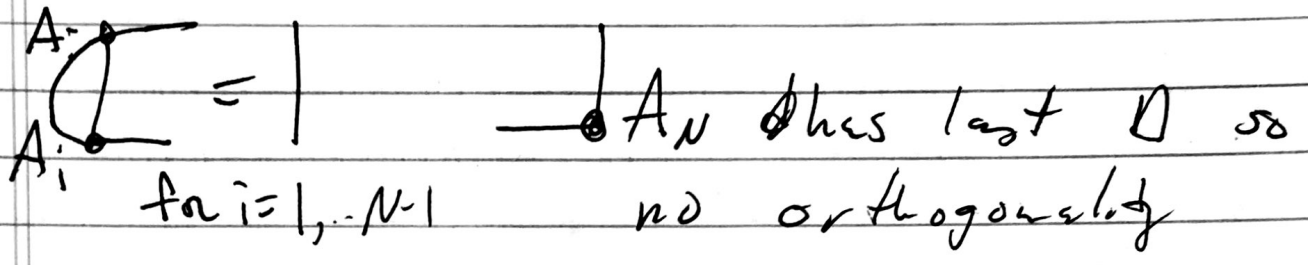


$$A_1 \begin{bmatrix} \\ \\ \end{bmatrix} = \begin{bmatrix} \\ \\ \end{bmatrix}$$

$$\begin{bmatrix} A_2 \\ A_2 \end{bmatrix} = \begin{bmatrix} \\ \end{bmatrix}$$

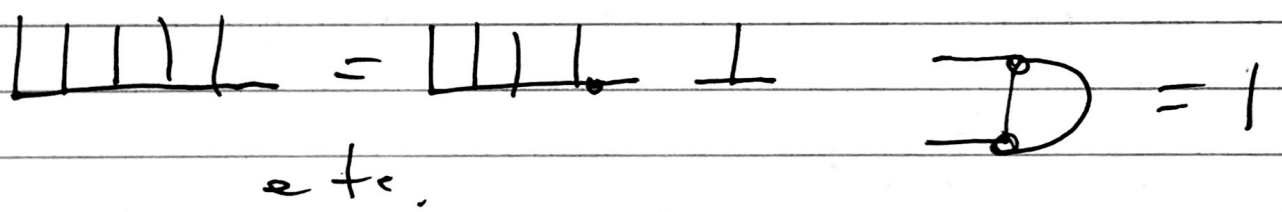
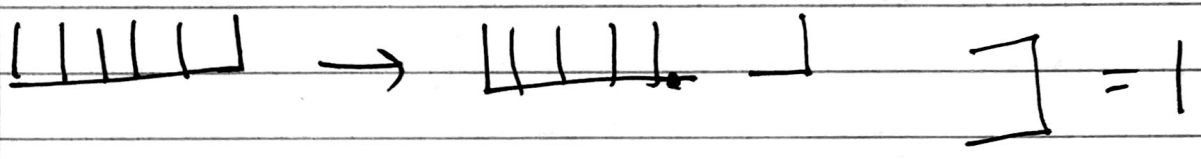


Keep going to the end

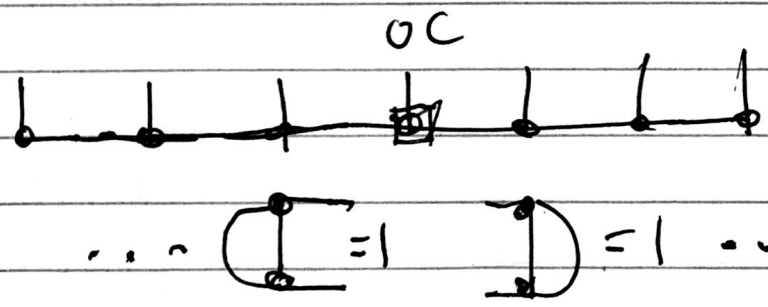


We say site N is the orthogonality center

We could put ~~it~~ ^{the OC} on ~~the~~ site 1

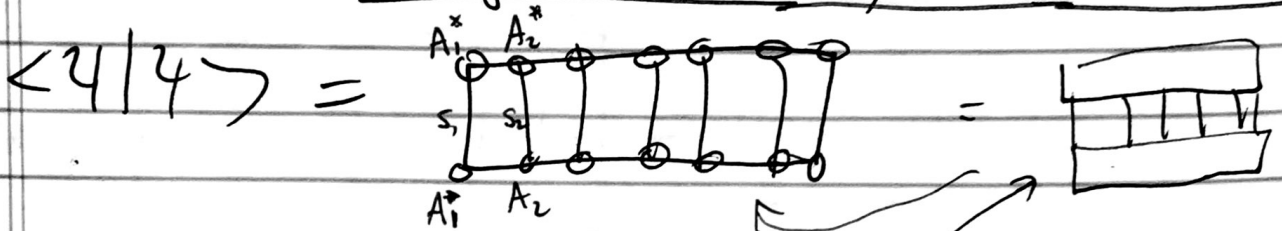


In general, it looks like

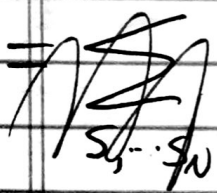


It is possible to not have an OC - but it is useful to have one.

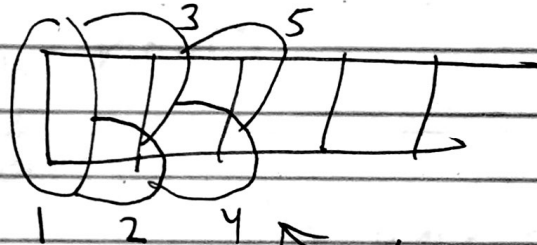
Normalization of an MPS, Measuring Ops



$$\sum_{S_1 \dots S_N} \sum_{S'_1 \dots S'_N} \psi_{S_1 \dots S_N} \psi_{S'_1 \dots S'_N} \langle S'_1 \dots S'_N | S_1 \dots S_N \rangle = \delta_{S_1 S'_1} \dots \delta_{S_N S'_N}$$



To contract: left to right, or right to left.



Not top, then bottom

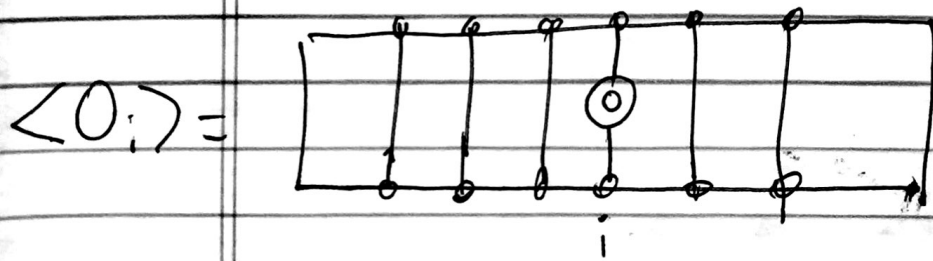
"Bubbling"

Measuring Operators

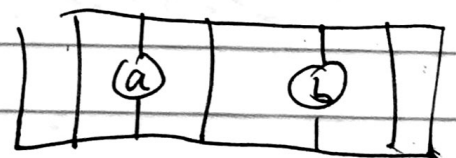
$$\langle S'_1 \dots S'_N | \hat{O}_i | S_1 \dots S_N \rangle$$

$$= \langle S'_1 \dots S'_N | S_1 \dots (O_i S_i) \dots S_N \rangle$$

O_i attaches to A_i



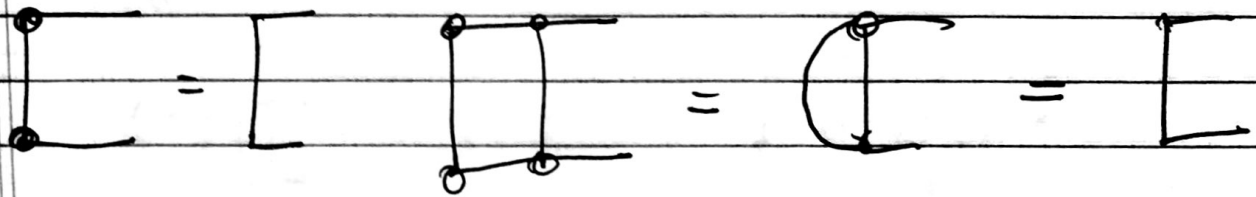
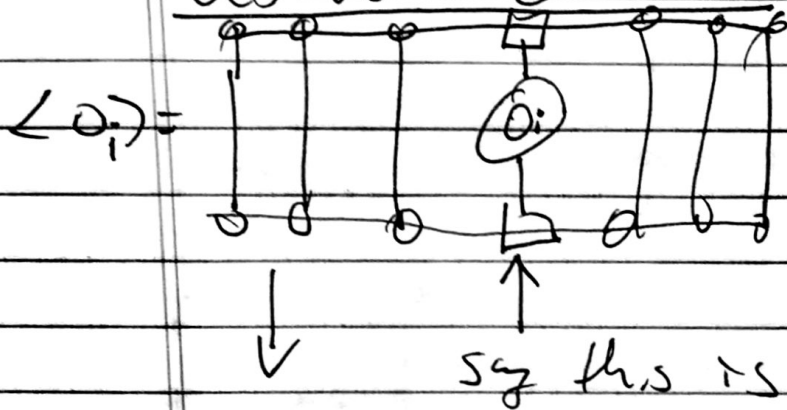
Also



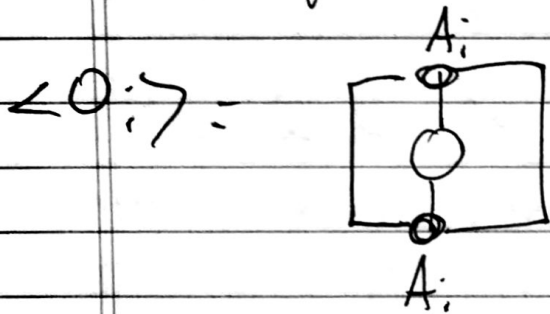
etc.

"Bubble" left to right,

Using an OC

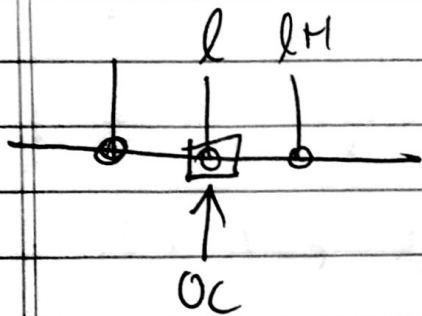


Similarly on right

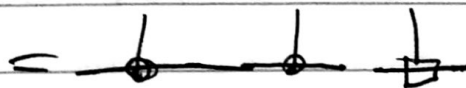
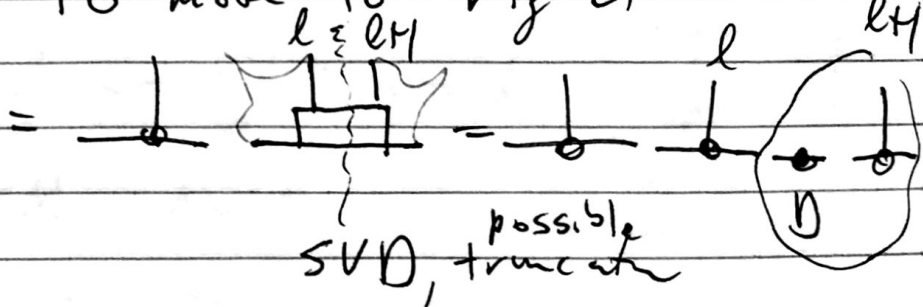


No need to do other contractions

Moving the OC

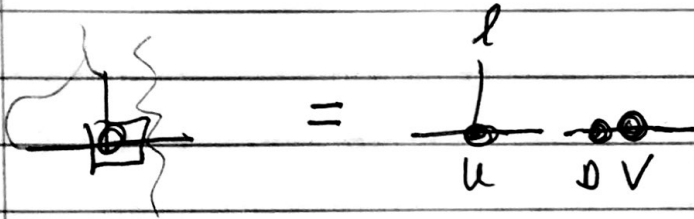


To move to right

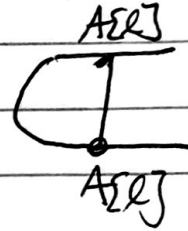
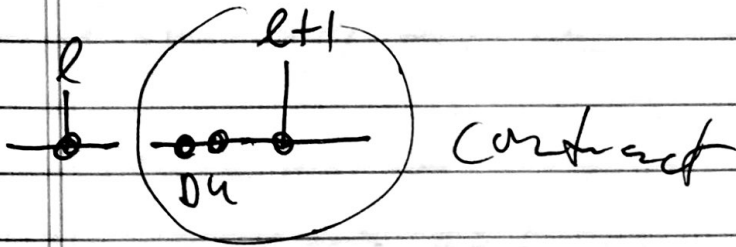


1) Join two adjacent tensors, then split, moving OC.

How: Using one site



then $A[l] = U$



$A[l+1] = \dots \text{ } DV \text{ } \{l+1\}$

One site, QR

The QR is a cheaper (than SVD) matrix factorization

where Q is

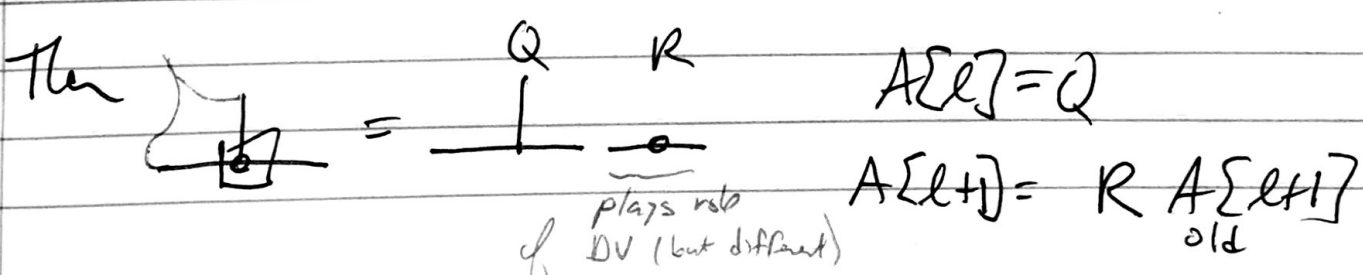
[row] orthogonal & R is

~~$M = QR$
 $N \times M \quad N \times M \quad M \times M$
 $N \geq M$~~

upper right triangular. The "thin" version

Make Q column-orthonormal

$$\begin{matrix} M \\ N \times M \end{matrix} = \begin{matrix} Q \\ N \times M \end{matrix} \begin{matrix} R \\ M \times M \end{matrix} \quad \begin{matrix} Q^T \\ Q \end{matrix} = \begin{matrix} I \\ 1 \end{matrix}$$



Getting the ground state ^{Directly} in MPS Form

~~TEBD / tDMRG~~ ~~(write here)~~
 time evolving block decomposition ~~April 28~~

Trotter Approximation:

Power method $V' = (1 - \epsilon H) V$ iterate
 $V \rightarrow$ ground state if ϵ small.

Better $V' = e^{-\beta H} V$ β could be big
 $e^{-\beta H}$ is hard for β big, but easy for
 β small, $\beta \rightarrow \tau$ $\tau \ll 1$

~~$e^{-\tau \sum_i \hat{S}_i \cdot \hat{S}_{i+1}}$~~ Let $H_i = \hat{S}_i \cdot \hat{S}_{i+1}$, $H = \sum_i H_i$
 $e^{-\tau H} = e^{-\tau \sum_i H_i} \approx \prod_i e^{-\tau H_i}$

Errors come from $[H_i, H_{i+1}] \neq 0$

$$e^{-\tau(A+B)} = 1 - \tau(A+B) + \frac{1}{2} \tau^2 \underbrace{(A+B)^2}_{A^2+B^2+AB+BA} + \dots$$

$$e^{-\tau A} e^{-\tau B} = \left(1 - \tau A + \frac{1}{2} \tau^2 A^2 \dots\right) \left(1 - \tau B + \frac{1}{2} \tau^2 B^2 \dots\right)$$

$$= 1 - \tau(A+B) + \frac{1}{2} \tau^2 (A^2 + B^2 + 2AB) \dots$$

The error is $\propto \tau^2$: $e^{-\tau(A+B)} = e^{-\tau A} e^{-\tau B} + O(\tau^2)$

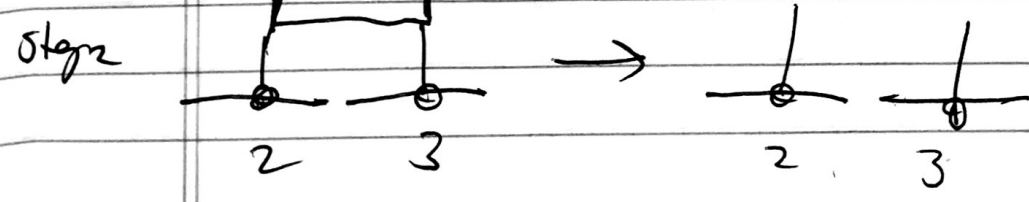
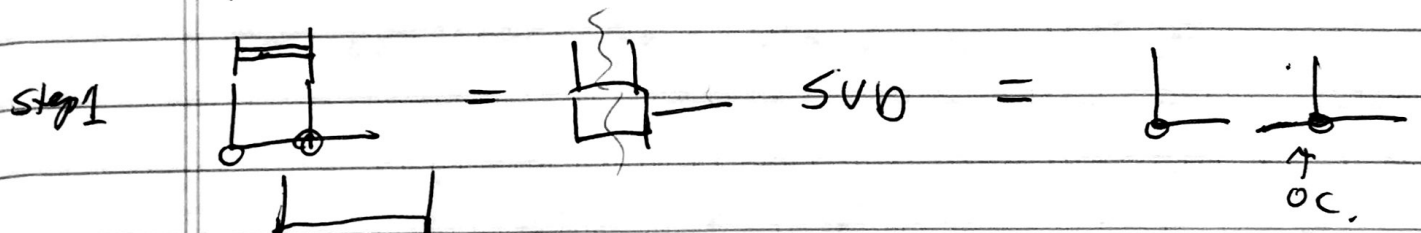
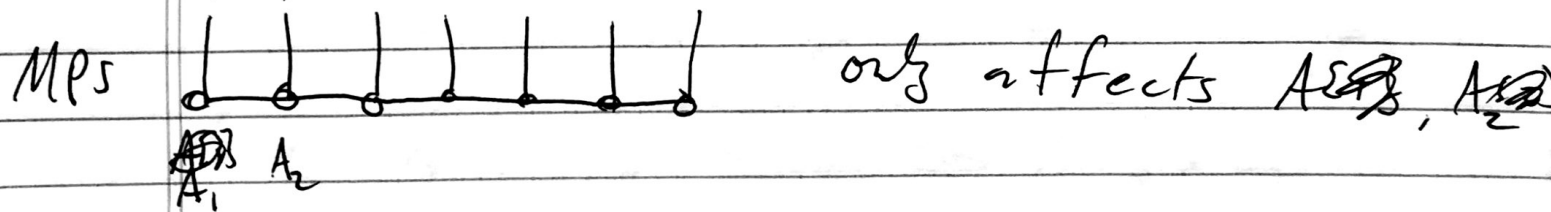
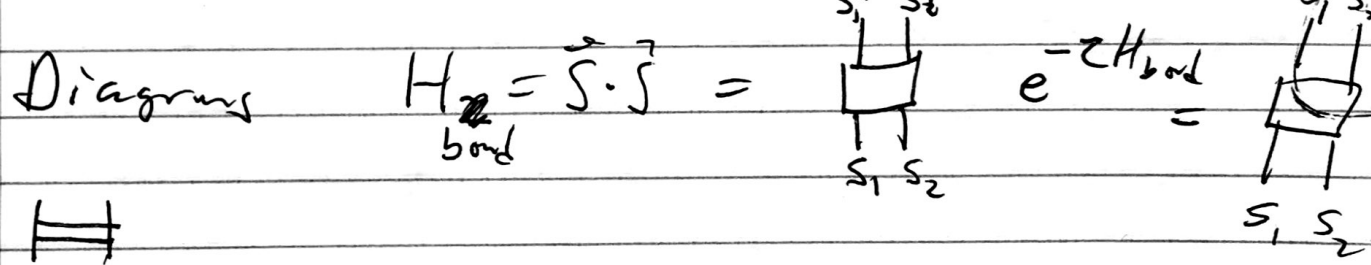
But to get good state, you need to do this $O(\frac{1}{\epsilon})$ times, so ~~total~~ total error $\sim \tau$ [Usually too lousy to use!]

If you do it twice, w. th a reverse, you get $O(\tau^3)$ per sweep $\rightarrow O(\tau^2)$ total

$$e^{-\tau(A+B+C)} \approx e^{-\tau A} e^{-\tau B} e^{-\tau C} e^{-\tau C} e^{-\tau B} e^{-\tau A}$$

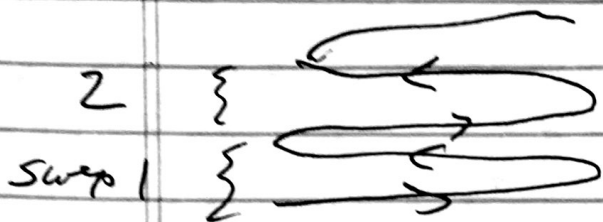
has AB+BA has CAB (no CAB) BA

We apply $e^{-\tau H_i}$ one at a time, moving the O_C in a sweeping algorithm



1 sweep = $i = 1:n$, then $i = n:-1:1$

Moving OC in direction of sweeping



Periodically, you need to fix the norm of the state.

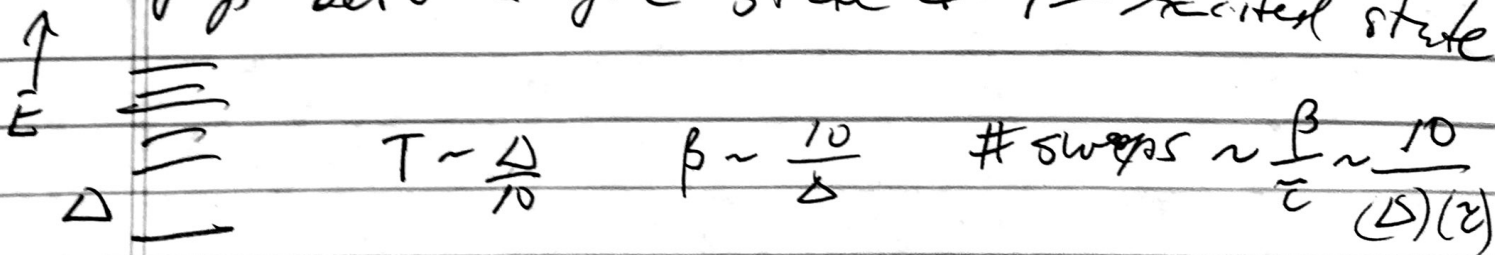
The splitting/moving of the OC needs the two side SVP; not QR.

QR is only when we don't truncate.

ϵ controls error in result

But needs to converge in # of sweeps

Need $T = \frac{1}{\epsilon}$ to be ~~less~~ less than gap between good state & 1st excited state



$$T \sim \frac{1}{\epsilon}$$

$$\beta \sim \frac{10}{\Delta}$$

$$\# \text{ sweeps} \sim \frac{\beta}{\epsilon} \sim \frac{10}{(\Delta)(\epsilon)}$$