\[ \mathcal{H} |\psi\rangle = E |\psi\rangle \]

Introduction to Exact Diagonalization

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Outline

\[ \mathcal{H} |\psi\rangle = E |\psi\rangle \]

- Main Idea
- Examples of typical applications
- Structure of an Exact Diagonalization package
  - Hilbertspace, Symmetries
  - Hamiltonian
  - Linear Algebra
  - Observables
  - Parallelization
Exact Diagonalization: Main Idea

- Solve the Schrödinger equation of a quantum many body system numerically

\[ \mathcal{H} |\psi\rangle = E |\psi\rangle \]

- Sparse matrix, but for quantum many body systems the vector space dimension grows exponentially!

- Some people will tell you that’s all there is.

- But if you want to get a maximum of physical information out of a finite system there is a lot more to do and the reward is a powerful:

  Quantum Mechanics Toolbox
Exact Diagonalization: Applications

- **Quantum Magnets**: nature of novel phases, critical points in 1D, dynamical correlation functions in 1D & 2D

- **Fermionic models (Hubbard/t-J)**: gaps, pairing properties, correlation exponents, etc

- **Fractional Quantum Hall states**: energy gaps, overlap with model states, entanglement spectra

- **Quantum dimer models or other constrained models (anyon chains, ...)**

- **Full Configuration Interaction in Quantum Chemistry**
Exact Diagonalization: Present Day Limits

- **Spin S=1/2 models:**
  - 40 spins square lattice, 39 sites triangular, 42 sites Honeycomb lattice
  - 64 spins or more in elevated magnetization sectors
  - Up to 1.5 billion (=10^9) basis states with symmetries, up to 4.5 billion without

- **t-J models:**
  - 32 sites checkerboard with 2 holes
  - 32 sites square lattice with 4 holes
  - Up to 2.8 billion basis states

- **Fractional quantum hall effect**
  - Different filling fractions $\nu$, up to 16-20 electrons
  - Up to 3.5 billion basis states

- **Hubbard models**
  - 20 sites square lattice at half filling, 21 sites triangular lattice
  - 22-25 sites in ultracold atoms setting w.o. spatial symmetries
  - Up to 160 billion basis states
Structure of an Exact Diagonalization code
Ingredients

- Hilbert space
  - Basis representation, Lookup techniques
  - Symmetries
- Hamiltonian Matrix
  - Sparse Matrix representation (memory/disk)
  - Matrix recalculation on the fly (matrix-free)
- Linear Algebra : Eigensolver / Time propagation
  - LAPACK full diagonalization
  - Lanczos type diagonalization (needs only $|v\rangle = H|u\rangle$ operations)
  - More exotic eigensolver techniques, real oder imaginary-time propagation,
- Observables
  - Static quantities (multipoint correlation functions, correlation density matrices,...)
  - Dynamic observables (spectral functions, density of states,...)
  - Real-time evolution
Hilbert Space
Basis representation

- States of the Hilbert space need to be represented in the computer.
- Choose a representation which makes it simple to act with the Hamiltonian or other operators on the states, and to localize a given state in the basis.
- Simple example: ensemble of S=1/2 sites in binary coding

\[
|\uparrow \uparrow \downarrow \uparrow\rangle \rightarrow [1 \ 1 \ 0 \ 1]_2 = 13
\]

Detection of up or down spin can be done with bit-test. Transverse exchange \(S^+ S^- + S^- S^+\) can be performed by an XOR operation:

\[
[1 \ 1 \ 0 \ 1]_2 \ XOR \ [0 \ 1 \ 1 \ 0]_2 = [1 \ 0 \ 1 \ 1]_2
\]

- For S=1, one bit is obviously not sufficient. Use ternary representation or simply occupy two bits to label the 3 states.
For t-J models at low doping it is useful to factorize hole positions and spin configurations on the occupied sites.

For Hubbard models one can factorize the Hilbert space in up and down electron configurations.

For constrained models - such as dimer models - the efficient generation of all basis states requires some thought.

One of the key challenges for a fast ED code is to find the index of the new configuration in the list of all configurations (index $f$ in $H_{i,f}$).

Let us look at the example of $S=1/2$ spins at fixed $S^Z$. 
One of the key problems for a fast ED code is to find the index of the new configuration in the list of all configurations (index $f$ in $H_{i,t}$).

$$[1\ 0\ 1\ 1]_2 = 11_{10}$$

But is 11 the index of this configuration in a list of all $S^z=1$ states? **no!**

Use Lin tables to map from binary number to index in list of allowed states:
(generalization of this idea works for arbitrary number of additive quantum numbers)

Two tables with $2^{(N/2)}$ [=sqrt($2^N$)] entries, one for MSBs and one for LSBs

<table>
<thead>
<tr>
<th>MSB</th>
<th>LSB</th>
<th>Ind([0 1 1 1])</th>
<th>Ind([1 0 1 1])</th>
<th>Ind([1 1 0 1])</th>
<th>Ind([1 1 1 0])</th>
</tr>
</thead>
<tbody>
<tr>
<td>[0 0] = X</td>
<td>[0 0] = X</td>
<td>0 + 0 = 0</td>
<td>1 + 0 = 1</td>
<td>2 + 0 = 2</td>
<td>2 + 1 = 3</td>
</tr>
</tbody>
</table>
Basis lookup procedures (Lin tables)

- Lookup can therefore be done with two direct memory reads. This is a time and memory efficient approach (at least in many interesting cases).

- An alternative procedure is to build a hash list [const access time] or to perform a binary search [log access time].

- This becomes somewhat more involved when using spatial symmetries...
Symmetries

Consider a XXZ spin model on a lattice. What are the symmetries of the problem?

\[ H = \sum_{i,j} J_{i,j}^{xy} (S_i^x S_j^x + S_i^y S_j^y) + J_{i,j}^{xz} S_i^z S_j^z \]

- The Hamiltonian conserves total \( S^z \), we can therefore work within a given \( S^z \) sector. This easily implemented while constructing the basis, as we discussed before.

- The Hamiltonian is invariant under the space group, typically a few hundred elements. (in many cases = Translations x Pointgroup). Needs some technology to implement...

- At the Heisenberg point, the total spin is also conserved. It is however very difficult to combine the SU(2) symmetry with the lattice symmetries in a computationally useful way (non-sparse and computationally expensive matrices).

- At \( S^z=0 \) one can use the spin-flip (particle-hole) symmetry which distinguishes even and odd spin sectors at the Heisenberg point. Simple to implement.
Spatial Symmetries

- Spatial symmetries are important for reduction of Hilbert space
- Symmetry resolved eigenstates teach us a lot about the physics at work, dispersion of excitations, symmetry breaking tendencies, topological degeneracy, ... ⇒ more about this in the second lecture

40 sites square lattice
\[ T \otimes PG = 40 \times 4 \text{ elements} \]

Icosidodecahedron (30 vertices)
\[ I_h: 120 \text{ elements} \]
Spatial Symmetries

- Symmetries are sometimes not easily visible, use graph theoretical tools to determine symmetry group [nauty, grape].
- In an ED code a spatial symmetry operation is a site permutation operation. (could become more complicated with spin-orbit interactions and multiorbital sites)

40 sites square lattice
T $\otimes$ PG = 40 x 4 elements

Icosidodecahedron (30 vertices)
I_h: 120 elements
Spatial Symmetries: Building the basis

- Build a list of all allowed states satisfying the “diagonal” constraints, like particle number, total $S^z$, ...

- For each state we apply all symmetry operations and keep the state as a representative if it has the smallest integer representation among all generated states in the orbit.

Example: 4 site ring with cyclic translation $T$, $S^z=1$ sector

<table>
<thead>
<tr>
<th>State</th>
<th>Operation</th>
<th>New State</th>
</tr>
</thead>
<tbody>
<tr>
<td>$T^0([0 1 1 1])$</td>
<td>$T^0([1 0 1 1])$</td>
<td>$[0 1 1 1]$</td>
</tr>
<tr>
<td>$T^1([0 1 1 1])$</td>
<td>$T^1([1 0 1 1])$</td>
<td>$[1 0 1 1]$</td>
</tr>
<tr>
<td>$T^2([0 1 1 1])$</td>
<td>$T^2([1 0 1 1])$</td>
<td>$[1 1 0 1]$</td>
</tr>
<tr>
<td>$T^3([0 1 1 1])$</td>
<td>$T^3([1 0 1 1])$</td>
<td>$[1 1 1 0]$</td>
</tr>
</tbody>
</table>

Keep state | Discard state
Spatial Symmetries: Building the basis

For one-dimensional representations $\chi$ of the spatial symmetry group:

- “Bloch” state
  $$ |\tilde{r}\rangle = \frac{1}{\mathcal{N} \sqrt{|G|}} \sum_{g \in G} \chi(g) |g(r)\rangle $$

- Norm of the state is given as:
  $$ \mathcal{N} = \sqrt{\sum_{g \in G, g(r)=r} \chi(g)} $$

- The norm (and therefore the state itself) can vanish if it has a nontrivial stabilizer combined with a nontrivial representation $\chi$.

Example: 4 site $S=1/2$ ring with cyclic translations:

<table>
<thead>
<tr>
<th>Site</th>
<th>$K=0$</th>
<th>$K=\pm\pi/2$</th>
<th>$K=\pi$</th>
<th>$2^4=16$</th>
</tr>
</thead>
<tbody>
<tr>
<td>$S^z=2$</td>
<td>$</td>
<td>1\ 1\ 1\ 1\rangle, \mathcal{N}=2$</td>
<td>$</td>
<td>0\ 1\ 1\ 1\rangle, \mathcal{N}=1$</td>
</tr>
<tr>
<td>$S^z=1$</td>
<td>$</td>
<td>0\ 1\ 1\ 1\rangle, \mathcal{N}=1$</td>
<td>$</td>
<td>0\ 1\ 1\ 1\rangle, \mathcal{N}=1$</td>
</tr>
<tr>
<td>$S^z=0$</td>
<td>$</td>
<td>0\ 0\ 1\ 1\rangle, \mathcal{N}=\sqrt{2}$</td>
<td>$</td>
<td>0\ 0\ 1\ 1\rangle, \mathcal{N}=1$</td>
</tr>
<tr>
<td></td>
<td></td>
<td></td>
<td></td>
<td>$4$</td>
</tr>
</tbody>
</table>
The Hamiltonian Matrix
The Hamiltonian Matrix

Now that we have a list of representatives and their norms, can we calculate the matrix elements of the Hamiltonian? \[ \langle \tilde{s} | H | \tilde{r} \rangle =? \]

Let us look at an elementary, non-branching term in the Hamiltonian:

\[ h^\alpha |r\rangle = h^\alpha (r) |s\rangle \]

We can now calculate the matrix element \[ \langle \tilde{s} | h^\alpha | \tilde{r} \rangle \] without double expanding the Bloch states:

\[ \langle \tilde{s} | h^\alpha | \tilde{r} \rangle = \frac{\mathcal{N}_s}{\mathcal{N}_r} \chi(g^*) h^\alpha (r) \]

key algorithmic problem: given a possibly non-representative \( |s\rangle \), how do we find the associated representative \( |\tilde{s}\rangle \), as well as a symmetry element \( g^* \) relating \( |s\rangle \) to \( |\tilde{s}\rangle \)?
The Hamiltonian Matrix

- key algorithmic problem: given a possibly non-representative \( |s\rangle \), how do we find the associated representative \( |\tilde{s}\rangle \), as well as a symmetry element \( g^* \) relating \( |s\rangle \) to \( |\tilde{s}\rangle \)?

- Brute force: loop over all symmetry operations applied on \( |s\rangle \) and retain \( |\tilde{s}\rangle \) and \( g^* \). This is however often not efficient (many hundred symmetries).

- Prepare a lookup list, relating each allowed configuration with the index of its representative, and also the associated group element linking the two. Gives fast implementation, but needs a list of the size of the non spatially-symmetrized Hilbert space.

- For specific lattices and models (Hubbard models) clever tricks exist which factorize the symmetry group into a sublattice conserving subgroup times a sublattice exchange. They give \( |\tilde{s}\rangle \) fast, then a hash or binary search is needed to locate \( |\tilde{s}\rangle \) in the list of representatives in order to get its index.
Hamiltonian Matrix Storage

Different possibilities exist:

- Store hamiltonian matrix elements in memory in a sparse matrix format. Fast matrix vector multiplies, but obviously limited by available memory.

- Store hamiltonian matrix elements on disk in a sparse matrix format. In principle possible due to the vast disk space available, but I/O speed is much slower than main memory access times. Difficult to parallelize.

- Recalculate the hamiltonian matrix elements in each iterations “on the fly”. Needed for the cutting edge simulations, where the whole memory is used by the Lanczos vectors. Can be parallelized on most architectures.
The Linear Algebra Backend
The Reference:

Lanczos Algorithm (C. Lanczos, 1950)

Three vector recursion

\[ |\phi'\rangle = H|\phi_n\rangle - \beta_n|\phi_{n-1}\rangle, \]
\[ \alpha_n = \langle \phi_n | \phi' \rangle, \]
\[ |\phi''\rangle = |\phi'\rangle - \alpha_n|\phi_n\rangle, \]
\[ \beta_{n+1} = ||\phi''|| = \sqrt{\langle \phi'' | \phi'' \rangle}, \]
\[ |\phi_{n+1}\rangle = |\phi''\rangle / \beta_{n+1}, \]

\[ \tilde{H}_N = \begin{bmatrix} \alpha_0 & \beta_1 & 0 & \cdots & 0 \\ \beta_1 & \alpha_1 & \beta_2 & 0 & \cdots \\ 0 & \beta_2 & \alpha_2 & \beta_3 & 0 \\ \vdots & \vdots & \vdots & \ddots & \ddots \\ 0 & \cdots & 0 & \beta_{N-2} & \alpha_{N-2} \\ 0 & \cdots & 0 & \beta_{N-1} & \alpha_{N-1} \end{bmatrix} \]

Eigenvalues of \( H_N \) converge rapidly towards eigenvalues of \( H \).

Once desired eigenvalue is converged, restart recursion and assemble the eigenvector.

very quick convergence for extremal eigenvalues!
Once the ground state has converged, the vectors in the recursion tend to lose their orthogonality. As a consequence fake new eigenvalues show up in the approximate spectrum. These can be removed by heuristic techniques.

Degeneracies of eigenvalues can not be resolved by construction. For this task one would need a band lanczos or the (Jacobi-)Davidson technique. However multiply degenerate eigenvalues are converged.

Checkpointing is useful when performing large-scale simulations.
Full Diagonalization: Thermodynamics

- Lapack / Householder complete diagonalization of the spectrum.

- Calculate partition function and all the thermodynamic quantities you want, often the only pedestrian method available for frustrated systems.

- Symmetries are also very important, because the computational requirements scale as $O(D^3)$, where D is the dimension of the block Hilbert space. Typical D’s for a workstation are a few 1’000, up to a few 100’000 on supercomputers.

Observables
Observables

- In principle once can calculate any correlation function, since one has access to the full many body wave functions. When using spatial symmetries, the correlation functions need to be properly symmetrized too.

- Complicated correlation functions occur in frustrated systems:

  \[
  \langle (\mathbf{S}_i \cdot \mathbf{S}_j)(\mathbf{S}_k \cdot \mathbf{S}_l) \rangle - \langle \mathbf{S}_i \cdot \mathbf{S}_j \rangle \langle \mathbf{S}_k \cdot \mathbf{S}_l \rangle
  \]

  \[
  \langle (\mathbf{S}_i \wedge \mathbf{S}_j)^z(\mathbf{S}_k \wedge \mathbf{S}_l)^z \rangle - \langle (\mathbf{S}_i \wedge \mathbf{S}_j)^z \rangle \langle (\mathbf{S}_k \wedge \mathbf{S}_l)^z \rangle
  \]

  Dimer-dimer correlations

  Spin current correlations
Frequency Dynamics

- $G_A(\omega + i\eta) = \langle \psi | A^\dagger \frac{1}{E_0 + \omega + i\eta - H} A | \psi \rangle$  
  $A = S^\alpha(q), c_k, \ldots$

- Generate Krylov space of $A|\psi\rangle$
  Use continued fraction used to invert $(E_0 + \omega + i\eta - H)$

- Triangular Lattice Spin Dynamics in zero field

AML unpublished
Exact Diagonalization
Real-Time Dynamics

It is expensive to obtain the full propagator $\exp[-i t H]$

Krylov methods exist to approximate the propagator for a given state $|\psi(0)\rangle$
One can get the time propagated state $|\psi(t)\rangle$ with only $|v\rangle = H|u\rangle$ operations.

Example: time evolution of a strongly correlated quantum systems after an abrupt change in the parameters in the Hamiltonian. Revivals and Relaxation.

C. Kollath, AML, E. Altman, PRL 2007
Parallelization Strategies
Parallelization: Shared memory nodes

- In the Lanczos algorithm the heaviest part is the elementary matrix-vector multiplication.
- In a matrix-free formulation this part can easily be parallelized using OpenMP pragmas in the code, even on your multi-core workstation. Choose the right strategy between pull and push!

In this parallelization we have uncritical concurrent reads, but no concurrent updates of vector v.
Parallelization: Shared memory nodes

- SGI Origin 3800 [Eridan]
- IBM p690 [zahir019/020]
- Sun
- IBM p5
- SGI Altix It2 1.5GHz

Time per 3 iterations [s]

- Scales well up to a few ten threads on “memory uniform” SMP machines.
Parallelization: Distributed memory nodes

- For some classes of problems the Hilbert space size is not too big, but the vast number of matrix elements is a challenge. [ED in momentum space formulation & Quantum Hall problems]

- These problems can be OpenMP parallelized, but are also suitable for large scale Message passing parallelization.
Parallelization: Distributed memory nodes

Strong scaling example RG-ED: matrix dimension 10 million performed on a 1024 node Cray XT-3 machine: speedup of $\approx 800$ on 1024 procs
Parallelization:
How to harness the petaflop computers?

- Cutting edge petaflop systems have a huge number of core, but only a moderate amount of node-local memory.

- Next generation ED codes need to be developed in order to attack e.g. the 80 billion Hilbert space of a 48 site kagome antiferromagnet.

- Problem remains difficult to parallelize due to its all-to-all structure. Global Arrays or UPC can help developing distributed ED codes.
Exact Diagonalization Literature

- N. Laflorencie & D. Poilblanc,
  "Simulations of pure and doped low-dimensional spin-1/2 gapped systems"

- R.M. Noack & S. Manmana,
  "Diagonalization- and Numerical Renormalization-Group-Based Methods for Interacting Quantum Systems",

- A. Weisse, H. Fehske
  "Exact Diagonalization Techniques"

- A. Läuchli
  "Numerical Simulations of Frustrated Systems"
Thank you!